

$$1. \quad p(0,0) = \frac{1}{3} \quad p(0,1) = \frac{1}{3} \quad p(1,0) = 0 \quad p(1,1) = \frac{1}{3}$$

$$p(x=0) = p(0,0) + p(0,1) = \frac{2}{3}$$

$$p(x=1) = p(1,0) + p(1,1) = \frac{1}{3}$$

$$p(y=0) = p(0,0) + p(1,0) = \frac{1}{3}$$

$$p(y=1) = p(0,1) + p(1,1) = \frac{2}{3}$$

$$a) \quad H(X) = -\frac{2}{3} \log \frac{2}{3} - \frac{1}{3} \log \frac{1}{3} = .528 + .390 = .918 \text{ bit}$$

$$b) \quad H(Y) = -\frac{1}{3} \log \frac{1}{3} - \frac{2}{3} \log \frac{2}{3} = .918 \text{ bit}$$

$$c) \quad H(X|Y) = -\sum_{i=1}^2 \sum_{j=1}^2 p(x_i, y_j) \log \frac{p(x_i, y_j)}{p(y_j)}$$

$$\begin{aligned} &= -\frac{1}{3} \log \frac{\frac{1}{3}}{\frac{1}{3}} - \frac{1}{3} \log \frac{\frac{1}{3}}{\frac{2}{3}} - 0 - \frac{1}{3} \log \frac{\frac{1}{3}}{\frac{2}{3}} \\ &= -\frac{2}{3} \log \frac{1}{2} = \frac{2}{3} \text{ bit} \end{aligned}$$

$$d) \quad H(XY) = H(Y) + H(X|Y)$$

$$= .918 + .667 = 1.585 \text{ bits}$$

$$e) I(X; Y) = H(X) - H(X|Y) = .918 - .667 = .251 \text{ bit}$$

$$f) H(Y|X) = - \sum_{i=1}^2 \sum_{j=1}^2 p(x_i, y_j) \log \frac{p(x_i, y_j)}{p(x_i)}$$

$$= -\frac{1}{3} \log \frac{1/3}{2/3} - \frac{1}{3} \log \frac{1/3}{1/3} - 0 - \frac{1}{3} \log \frac{1/3}{2/3}$$

$$= -\frac{2}{3} \log \frac{1}{2} = \frac{2}{3} \text{ bit}$$

2) source X with $N = 8$ symbols

the smallest entropy occurs when X is deterministic

$$p(x_i) = 1 \text{ for some } i$$

$$H(X) = 0 \text{ bit}$$

the largest entropy occurs when the symbols are equiprobable

$$p(x_i) = 1/8 \text{ for all } i$$

$$H(X) = \log_2 8 = 3 \text{ bits}$$

$$3. H(X) = H(X|Y) + I(X;Y)$$

$$H(Y) = H(Y|X) + I(X;Y)$$

$$H(XY) = H(X|Y) + H(Y|X) + I(X;Y)$$

$$\text{o.o} \quad H(X) + H(Y) - H(XY) = I(X;Y) \geq 0$$

Equality holds when $I(X;Y) = 0$

this occurs when X and Y are statistically independent

4. the relative entropy is minimized when

$$q(x) = p(x)$$

therefore $q(x_1) = 1/2 \quad q(x_2) = 1/4$

$$q(x_3) = q(x_4) = 1/8$$